## Real Analysis: Solutions to Exercise Sheet 2

1. (a) Suppose, for a contradiction, that a > b, then a - b > 0. As  $a < b + \epsilon$  for all  $\epsilon > 0$ , this should be true for  $\epsilon = a - b$ . So we must have

$$a < b + (a - b) = a$$

But this is a contradiction (cannot have a < a). Thus our assumption was false and we must have  $a \leq b$  as required.

(b) Start with ' $\Rightarrow$ '. We assume that  $|x-a| < \epsilon$  and we want to prove that  $a-\epsilon < x < a+\epsilon$ . Either  $x-a \ge 0$  or x-a < 0, we will consider these two cases separately. First suppose  $x-a \ge 0$ , then

$$|x - a| = |x - a| < \epsilon$$

and so adding a on both sides using (O2) we get  $x < a + \epsilon$ . On the other hand, as  $x - a \ge 0$  and  $\epsilon > 0$  we have  $x - a > -\epsilon$ . Adding a on both sides using (O2) we get  $x > a - \epsilon$ . So we are done in this case. Next suppose x - a < 0, then

$$-(x-a) = |x-a| < \epsilon$$

and multiplying both sides by (-1) using (O4) we get  $(x-a) > -\epsilon$ . Now adding a on both sides using (O2) we get  $x > a - \epsilon$ . On the other hand x - a < 0 and  $\epsilon > 0$  so we have  $x - a < \epsilon$ . Adding a on both sides we get  $x < a + \epsilon$ .

Now we turn to ' $\Leftarrow$ '. We assume that  $a - \epsilon < x < a + \epsilon$  and we want to prove that  $|x - a| < \epsilon$ . Subtracting a on both sides (using (O2) with c = -a), we get

$$-\epsilon < x - a < \epsilon.$$

If  $x - a \ge 0$  then  $|x - a| = x - a < \epsilon$  so we are done. If x - a < 0 then |x - a| = -(x - a). Now as  $-\epsilon < x - a$  we get  $\epsilon > -(x - a)$  (using (O4) with c = -1). Thus we get  $|x - a| = -(x - a) < \epsilon$  as required.

- (c) For all  $\epsilon > 0$  we have  $|x a| < \epsilon$ . Using (b) this implies that  $x < a + \epsilon$  for all  $\epsilon > 0$ . But then using (a), we have that  $x \le a$ . On the other hand, we also get from (b) that  $a < x + \epsilon$  for all  $\epsilon > 0$ . So using (a) we get that  $a \le x$ . We have proved that at the same time  $a \le x$  and  $x \le a$ , so we must have x = a.
- 2. (a) This set is equal to S = [a, b]. max  $S = b = \sup S$ , min  $S = a = \inf S$ .
  - (b) This set is equal to S = (a, b]. max  $S = b = \sup S$ , S has no minimum, inf S = a.
  - (c) This set is equal to S = [a, b). S has no maximum,  $\sup S = b$ ,  $\min S = a = \inf S$ .

- (d)  $\min S = -3 = \inf S$ ,  $\max S = 103 = \sup S$ .
- (e)  $\min S = 0 = \inf S$ ,  $\max S = \frac{3}{2} = \sup S$ .
- (f) This set is not bounded above, not bounded below, it has no maximum, no minimum, no supremum and no infimum.
- (g) The roots of the polynomial  $3x^2 10x + 3$  are given by  $\frac{1}{3}$  and 3 and the graph of the function  $f(x) = 3x^2 10x + 3$  is 'facing upward'. Thus  $3x^2 10x + 3 < 0$  precisely when  $\frac{1}{3} < x < 3$ . Thus this set can be rewritten as  $(\frac{1}{3}, 3)$ . So it has no minimum, no maximum, its infimum is  $\frac{1}{3}$  and its supremum is 3.
- (h) This set has no maximum, no minimum, its supremum is  $\frac{1}{2}$  and its infimum is  $-\frac{1}{3}$ .
- (i) Using Exercise 1(b) above, this set can be rewritten as

$$\{x \in \mathbb{R} : -1 < x < 3\} = (-1, 3),$$

so no maximum, no minimum, infimum equal to -1 and supremum equal to 3.

3. Let  $x \in (0, \infty)$ , then x > 0. We show that  $y = \frac{x}{2}$  has the right property. Starting with the inequality x > 0 and using (O3) (with  $c = \frac{1}{2}$ ), we get  $\frac{x}{2} > 0$ . Now adding  $\frac{x}{2}$  on both sides (using (O2)) we get  $x > \frac{x}{2}$ . Thus  $0 < \frac{x}{2} < x$  and we can take  $y = \frac{x}{2} \in (0, \infty)$ . Now we need to deduce that  $(0, \infty)$  has no minimum. Suppose, for a contradiction,

that it had a minimum  $m \in (0, \infty)$ . So m > 0, but by the above argument, we can find  $y \in (0, \infty)$  (namely  $y = \frac{m}{2}$ ) such that y < m. This contradicts the fact that m is a minimum. Thus  $(0, \infty)$  has no minimum.

4. First we show that 1 is an upper bound. We have -1 < 0 and adding n on both sides using (O2) we get n - 1 < n. Now dividing by n on both sides (i.e. using (O3) with  $c = \frac{1}{n} > 0$ ) we get  $\frac{n-1}{n} < 1$ . This is true for all  $n \in \mathbb{N}$ , so 1 is an upper bound for S.

We still need to show that 1 is the smallest upper bound. Suppose, for a contradiction, that it is not the smallest upper bound. Then S has a smaller upper bound  $\alpha$ , say, with  $\alpha < 1$ . Using Corollary 1.3.2 from the lecture we can find a rational number between  $\alpha$  and 1

$$\alpha < \frac{p}{q} < 1.$$

As  $\frac{p}{q} < 1$  we have  $p \leq q - 1$  and so

$$\alpha < \frac{p}{q} \le \frac{q-1}{q}.$$

But this means that  $\alpha$  is not an upper bound for S as if we take n = q we have  $\frac{n-1}{n} > \alpha$ . This is a contradiction and so 1 is the smallest upper bound.